# Math 530: Graph Theory, Spring 2022: Homework 8 due 2022-06-08 at 11:00 AM

## Please solve 4 of the 8 problems!

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## 1 Exercise 1

#### 1.1 Problem

Let  $D = (V, A, \psi)$  be a strongly connected multidigraph.

A wealth distribution on D shall mean a family  $(k_v)_{v \in V}$  of integers (one for each vertex  $v \in V$ ). If  $k = (k_v)_{v \in V}$  is a wealth distribution, then we refer to each value  $k_v$  as the wealth of the vertex v, and we define the total wealth of k to be the sum  $\sum_{v \in V} k_v$ . We say that a vertex v is in debt in a given wealth distribution  $k = (k_v)_{v \in V}$  if its wealth  $k_v$  is negative.

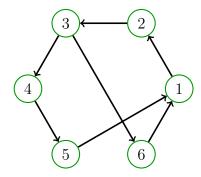
For any vertices v and w, we let  $a_{v,w}$  denote the number of arcs that have source v and w.

A donation is an operation that transforms a wealth distribution as follows: We choose a vertex v, and we decrease its wealth by its outdegree  $\deg^+ v$ , and then increase the wealth of each vertex  $w \in V$  (including v itself) by  $a_{v,w}$ . (You can think of v as donating a unit of wealth for each arc that has source v. This unit flows to the target to this arc. Note that a donation does not change the total wealth.)

Let k be a wealth distribution on D whose total wealth is larger than |A| - |V|. Prove that by an appropriately chosen finite sequence of donations, we can ensure that no vertex is in debt.

## 1.2 Remark

For instance, consider the digraph



with wealth distribution  $(k_1, k_2, k_3, k_4, k_5, k_6) = (-1, -1, 1, 2, 0, 1)$ . The vertices 1 and 2 are in debt here, but it is possible to get all vertices out of debt by having the vertices 4, 5, 6, 1 donate in some order (the order clearly does not matter for the result<sup>1</sup>).

Note that vertices are allowed to donate multiple times (although in the above example, this was unnecessary).

#### 1.3 HINT

Show first that if the total wealth is larger than |A| - |V|, then at least one vertex v has wealth  $\geq \deg^+ v$ .

#### 1.4 SOLUTION

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## 2 Exercise 2

#### 2.1 Problem

We continue with the setting and terminology of Exercise 1.

A clawback is an operation that transforms a wealth distribution as follows: We choose a vertex v, and we increase its wealth by its outdegree  $\deg^+ v$ , and then decrease the wealth of each vertex  $w \in V$  (including v itself) by  $a_{v,w}$ . (Thus, a clawback is the inverse of a donation.)

Let k be a wealth distribution on D whose total wealth is larger than |A| - |V|. Prove that by an appropriately chosen finite sequence of clawbacks, we can ensure that no vertex is in debt.

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<sup>&</sup>lt;sup>1</sup>Depending on the order, some vertices will go into debt in the process, but this is okay as long as they ultimately end up debt-free.

#### 2.2 Remark

Note that we are still assuming D to be strongly connected. Otherwise, the truth of the claim is not guaranteed. For instance, for the digraph



with wealth distribution  $(k_1, k_2, k_3, k_4) = (0, 0, -1, 2)$ , no sequence of donations and claw-backs will result in every vertex being out of debt (since the wealth difference  $k_4 - k_3$  is preserved under any donation or clawback, but this difference is too large to come from a debt-free distribution with total weight 1).

#### 2.3 HINT

Show that any donation is equivalent to an appropriately chosen composition of clawbacks.

#### 2.4 SOLUTION

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## 3 Exercise 3

#### 3.1 Problem

Let G = (V, E) be a simple graph such that each vertex of G has degree  $\geq 1$ . Prove that there exists a subset S of V having size  $\geq \sum_{v \in V} \frac{2}{1 + \deg v}$  and with the property that the induced subgraph G[S] is a forest.

#### 3.2 Remark

As the example of 2 3 shows, this claim is not true for loopless multigraphs (unlike the similar Theorem 1.1.3 in Lecture 23).

#### 3.3 SOLUTION

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## 4 Exercise 4

#### 4.1 Problem

Let n be a positive integer. Prove that there exists a tournament with n vertices and at least  $\frac{n!}{2^{n-1}}$  Hamiltonian paths.

## 4.2 SOLUTION

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## 5 Exercise 5

#### 5.1 Problem

Let  $G = (V, E, \varphi)$  be a multigraph. Let M be a matching of G.

An augmenting path for M shall mean a path  $(v_0, e_1, v_1, e_2, v_2, \dots, e_k, v_k)$  of G such that k is odd<sup>2</sup> and such that

- the even-indexed edges  $e_2, e_4, \ldots, e_{k-1}$  belong to M (note that this condition is vacuously true if k = 1);
- the odd-indexed edges  $e_1, e_3, \ldots, e_k$  belong to  $E \setminus M$ ;
- neither the starting point  $v_0$  nor the ending point  $v_k$  is matched in M.

Prove that M has maximum size among all matchings of G if and only if there exists no augmenting path for M.

#### 5.2 Hint

If M and M' are two matchings of G, what can you say about the symmetric difference  $(M \cup M') \setminus (M \cap M')$ ?

#### 5.3 SOLUTION

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<sup>&</sup>lt;sup>2</sup>Note that k = 1 is allowed.

## 6 Exercise 6

#### 6.1 Problem

Let (G, X, Y) be a bipartite graph. Prove that

$$\sum_{A \subseteq X} (-1)^{|A|} [N(A) = Y] = \sum_{B \subseteq Y} (-1)^{|B|} [N(B) = X]$$

(where we are using the Iverson bracket notation).

6.2 SOLUTION

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## 7 Exercise 7

#### 7.1 Problem

Let A and B be two finite sets such that  $|B| \ge |A|$ . Let  $d_{i,j}$  be a real number for each  $(i,j) \in A \times B$ . Let<sup>3</sup>

$$m_1 = \min_{\sigma: A \to B \text{ injective}} \max_{i \in A} d_{i,\sigma(i)}$$

and

$$m_2 = \max_{\substack{I \subseteq A; \ J \subseteq B; \\ |I| + |J| = |B| + 1}} \min_{(i,j) \in I \times J} d_{i,j}.$$

Prove that  $m_1 = m_2$ .

7.2 SOLUTION

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### 8 Exercise 8

#### 8.1 Problem

Let p be a prime number. Let  $(a_1, a_2, a_3, \ldots)$  be a sequence of integers that is periodic with period p (that is, that satisfies  $a_i = a_{i+p}$  for each i > 0). Assume that  $a_1 + a_2 + \cdots + a_p$  is not divisible by p. Prove that there exists an  $i \in \{1, 2, \ldots, p\}$  such that none of the p numbers

$$a_i, a_i + a_{i+1}, a_i + a_{i+1} + a_{i+2}, \ldots, a_i + a_{i+1} + \cdots + a_{i+p-1}$$

(that is, of the p sums  $a_i + a_{i+1} + \cdots + a_j$  for  $i \leq j < i + p$ ) is divisible by p.

<sup>&</sup>lt;sup>3</sup>The notation "min<sub>some kind of objects</sub> some kind of value" means the minimum of the given value over all objects of the given kind. An analogous notation is used for a maximum.

## 8.2 Remark

This would be false if p was not prime. For instance, for p=4, the sequence  $(0,2,2,2,0,2,2,2,\ldots)$  would be a counterexample.

## 8.3 HINT

What is the digraph, and why does it have a cycle?

8.4 SOLUTION

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