# Math 530: Graph Theory, Spring 2022: Homework 3 due 2022-04-20 at 11:00 AM Please solve 3 of the 6 problems!

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# 1 EXERCISE 1

## 1.1 PROBLEM

Which of the exercises 1, 2, 3, 4, 6 from homework set #1 remain true if "simple graph" is replaced by "multigraph"?

(For each exercise that becomes false, provide a counterexample. For each exercise that remains true, either provide a new solution that works for multigraphs, or argue that the solution we have seen applies verbatim to multigraphs, or derive the multigraph case from the simple graph case.)

## 1.2 Solution

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## 2 EXERCISE 2

#### 2.1 PROBLEM

Let G be a multigraph with at least one vertex. Let d > 2 be an integer. Assume that  $\deg v > 2$  for each vertex v of G. Prove that G has a cycle whose length is not divisible by d.

#### 2.2 Solution

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# 3 EXERCISE 3

#### 3.1 Problem

Let  $G = (V, E, \varphi)$  be a multigraph that has no loops.

If  $e \in E$  is an edge that contains a vertex  $v \in V$ , then we let e/v denote the endpoint of e distinct from v. (If e is a loop, then this is understood to mean v itself.)

For each  $v \in V$ , we define a rational number  $q_v$  by

$$q_v = \sum_{\substack{e \in E; \\ v \in \varphi(e)}} \frac{\deg(e/v)}{\deg v}.$$

(Note that the denominator  $\deg v$  on the right hand side is nonzero whenever the sum is nonempty!)

(Thus,  $q_v$  is the average degree of the neighbors of v, weighted with the number of edges that join v to the respective neighbors. If v has no neighbors, then  $q_v = 0$ .)

Prove that

$$\sum_{v \in V} q_v \ge \sum_{v \in V} \deg v.$$

(In other words, in a social network, your average friend has, on average, more friends than you do!)

#### 3.2 Hint

Any two positive reals x and y satisfy  $\frac{x}{y} + \frac{y}{x} \ge 2$ . Why, and how does this help?

#### 3.3 SOLUTION

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# 4 EXERCISE 4

## 4.1 PROBLEM

Let n and k be two integers such that n > k > 0. Define the simple graph  $Q_{n,k}$  as follows: Its vertices are the bitstrings  $(a_1, a_2, \ldots, a_n) \in \{0, 1\}^n$ ; two such bitstrings are adjacent if and only if they differ in exactly k bits<sup>1</sup>. (Thus,  $Q_{n,1}$  is the n-hypercube graph  $Q_n$ .)

- (a) Does  $Q_{n,k}$  have a hamc<sup>2</sup> when k is even?
- (b) Does  $Q_{n,k}$  have a hamc when k is odd?

## 4.2 Remark

One way to approach part (b) is by identifying the set  $\{0, 1\}$  with the field  $\mathbb{F}_2$  with two elements. The bitstrings  $(a_1, a_2, \ldots, a_n) \in \{0, 1\}^n$  thus become the size-*n* row vectors in the  $\mathbb{F}_2$ -vector space  $\mathbb{F}_2^n$ . Let  $e_1, e_2, \ldots, e_n$  be the standard basis vectors of  $\mathbb{F}_2^n$  (so that  $e_i$  has a 1 in its *i*-th position and zeroes everywhere else). Then, two vectors are adjacent in the *n*-hypercube graph  $Q_n$  (resp. in the graph  $Q_{n,k}$ ) if and only if their difference is one of the standard basis vectors (resp., a sum of k distinct standard basis vectors). Try to use this to find a graph isomorphism from  $Q_n$  to a subgraph of  $Q_{n,k}$ .

## 4.3 Solution

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# 5 EXERCISE 5

## 5.1 Problem

Let  $G = (V, E, \varphi)$  be a connected multigraph with 2m edges, where  $m \in \mathbb{N}$ . A set  $\{e, f\}$  of two distinct edges will be called a *friendly couple* if e and f have at least one endpoint in common. Prove that the edge set of G can be decomposed into m disjoint<sup>3</sup> friendly couples (i.e., there exist m disjoint friendly couples  $\{e_1, f_1\}, \{e_2, f_2\}, \ldots, \{e_m, f_m\}$  such that  $E = \{e_1, f_1, e_2, f_2, \ldots, e_m, f_m\}$ ).

**[Example:** Here is a graph with an even number of edges:



<sup>&</sup>lt;sup>1</sup>In other words: Two vertices  $(a_1, a_2, \ldots, a_n)$  and  $(b_1, b_2, \ldots, b_n)$  are adjacent if and only if the number of  $i \in \{1, 2, \ldots, n\}$  satisfying  $a_i \neq b_i$  equals k.

<sup>&</sup>lt;sup>2</sup>Recall that "hamc" is short for "Hamiltonian cycle".

<sup>&</sup>lt;sup>3</sup>"Disjoint" means "disjoint as sets" – i.e., having no edges in common.

One possible decomposition into disjoint friendly pairs is  $\{a, y\}, \{b, z\}, \{c, x\}$ .]

#### 5.2 HINT

Induct on |E|. Pick a vertex v of degree > 1 and consider the components of  $G \setminus v$ .

#### 5.3 Solution

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# 6 EXERCISE 6

### 6.1 PROBLEM

Let  $n \ge 0$ . Let  $d_1, d_2, \ldots, d_n$  be n nonnegative integers such that  $d_1 + d_2 + \cdots + d_n$  is even.

- (a) Prove that there exists a multigraph G with vertex set  $\{1, 2, ..., n\}$  such that all  $i \in \{1, 2, ..., n\}$  satisfy deg  $i = d_i$ .
- (b) A multigraph is said to be *loopless* if it has no loops. Prove that there exists a loopless multigraph G with vertex set  $\{1, 2, ..., n\}$  such that all  $i \in \{1, 2, ..., n\}$  satisfy deg  $i = d_i$  if and only if each  $i \in \{1, 2, ..., n\}$  satisfies the inequality

$$\sum_{\substack{j \in \{1,2,\dots,n\};\\j \neq i}} d_j \ge d_i.$$

$$\tag{1}$$

## 6.2 Remark

The inequality (1) is the "*n*-gon inequality": It is equivalent to the existence of a (possibly degenerate) *n*-gon with sidelengths  $d_1, d_2, \ldots, d_n$ .

#### 6.3 SOLUTION

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