

The twelvefold way so far:

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f is...	arbitrary	injective	surjective
$A \rightarrow X$			
$L \rightarrow L$	$ X ^{ A }$	$ X (X -1)(X -2) \dots (X - A +1)$	$\text{sur}(A , X) \leftarrow = X ! \binom{ A }{ X }$
$U \rightarrow L$	$\binom{ A + X -1}{ A }$	$\binom{ X }{ A }$	$\binom{ A -1}{ A - X }$
$L \rightarrow U$	$\binom{ A }{0} + \binom{ A }{1} + \dots + \binom{ A }{ X }$	$[A \leq X]$	$\binom{ A }{ X }$
$U \rightarrow U$			

More on Stirling numbers of 2nd kind:

(Recall: If $n \in \mathbb{N}$ and $k \in \mathbb{N}$, then $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \frac{1}{k!} \text{sur}(n, k)$
is the # of all set partitions of $[n]$ into k parts.)

Prop. 4.11. Let $n \in \mathbb{N}$ and $k \in \mathbb{N}$,

(a) We have $\left\{ \begin{matrix} n \\ 0 \end{matrix} \right\} = [n=0]$,

(b) $\left\{ \begin{matrix} 0 \\ k \end{matrix} \right\} = [k=0]$,

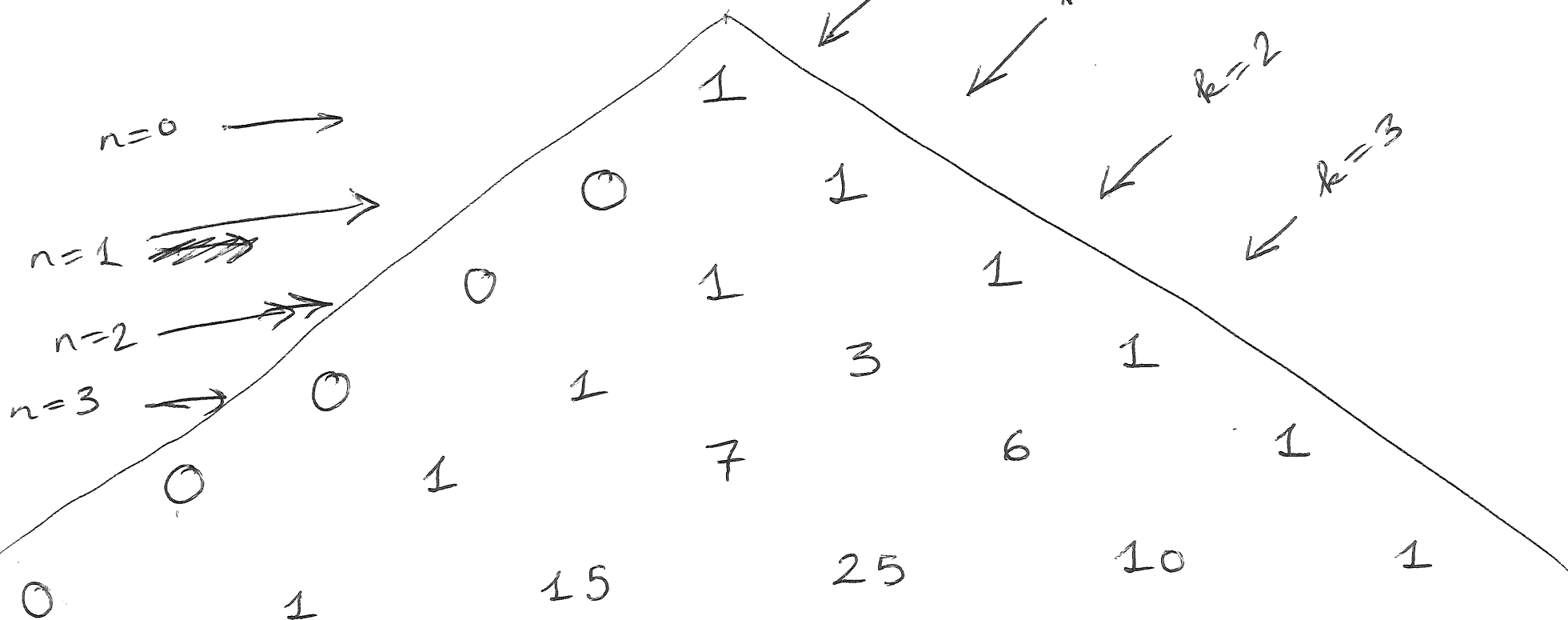
(c) $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = 0$ if $k > n$,

(d) $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \left\{ \begin{matrix} n-1 \\ k-1 \end{matrix} \right\} + k \left\{ \begin{matrix} n-1 \\ k \end{matrix} \right\}$ if $n > 0$ & $k > 0$,

(e) $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \sum_{j=0}^{n-1} \binom{n}{j} \left\{ \begin{matrix} j \\ k-1 \end{matrix} \right\} / k$.

Proof. Follows from Prop. 3.10, Prop. 3.11, Prop. 3.12. \square

Pascal-like triangle for $\binom{n}{k}$:

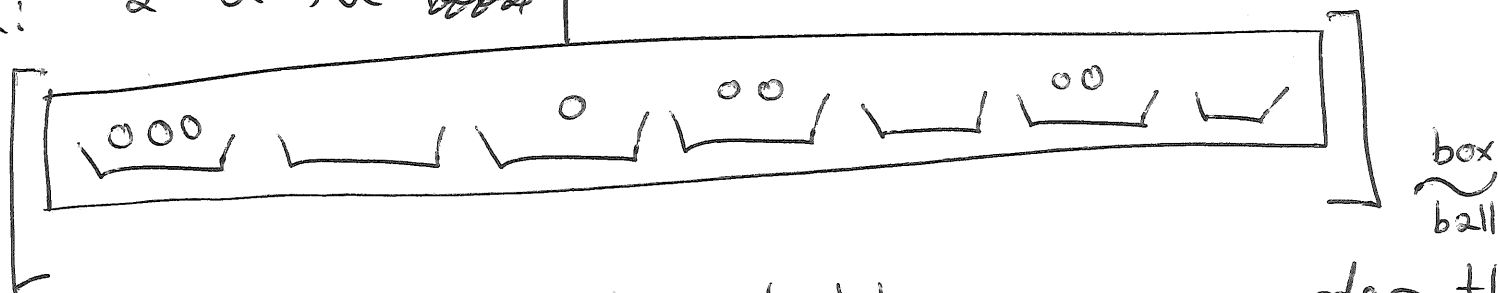


Note: If p is prime, then $p \mid \binom{p}{k} \forall k \in \{2, 3, \dots, p-1\}$.
 Can you prove it? (Compare to $p \mid \binom{p}{k} \forall k \in \{1, 2, \dots, p-1\}$, but this is harder.)

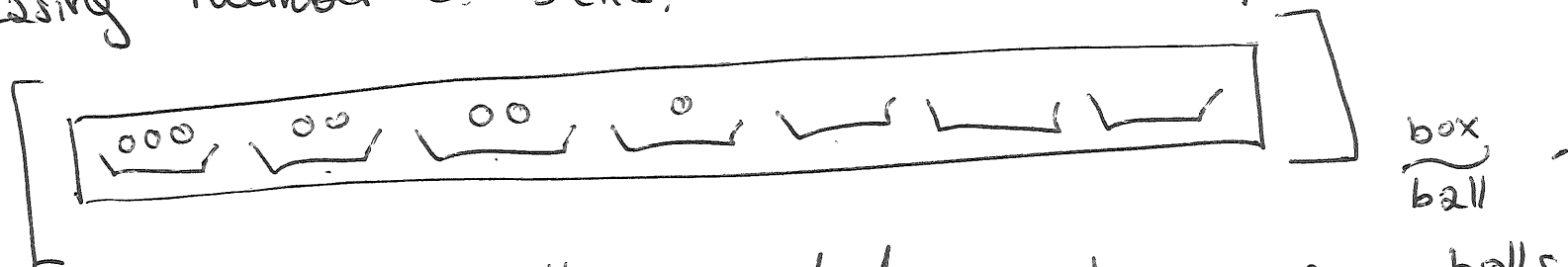
4.6. $U \rightarrow U$, AND INTEGER PARTITIONS

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Idea: a $U \rightarrow U$ ~~box~~ placement has the form



Since the boxes are indistinguishable, you can order them by decreasing number of balls. So the above placement is



Such a placement is thus encoded by how many balls each box has: here, the numbers are 3, 2, 2, 1, 0, 0, 0.

The decreasing order makes this encoding unique.

Def. A partition of an integer n is a weakly decreasing list (a_1, a_2, \dots, a_k) of positive integers whose sum is n (that is, $a_1 \geq a_2 \geq \dots \geq a_k > 0$ and $a_1 + a_2 + \dots + a_k = n$).

The integers a_1, a_2, \dots, a_k are called the parts of the partition.

Example: The partitions of 5 are
(5), (4, 1), (3, 2), (3, 1, 1), (2, 2, 1),
(2, 1, 1, 1), (1, 1, 1, 1, 1).

The only partition of 0 is ().

~~Def~~ (So a partition is a weakly decreasing composition.)

Def. Let $n \in \mathbb{Z}$ and $k \in \mathbb{N}$. Then, $p_k(n)$ means the # of partitions of n into k parts (=having k parts).

Example: $p_0(5) = 0$, $p_1(5) = 1$, $p_2(5) = 2$, $p_3(5) = 2$,
 $p_4(5) = 1$, $p_5(5) = 1$, $p_k(5) = 0 \forall k > 5$.

Prop. 4.12. (a) $p_k(n) = 0 \quad \forall n < 0,$

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(b) $p_k(n) = 0$ if $k > n.$

(c) $p_0(n) = [n=0].$

(d) $p_1(n) = [n \neq 0].$

(e) $p_k(n) = p_k(n-k) + p_{k-1}(n-1) \quad \forall k \geq 1 \text{ and } n \in \mathbb{Z}.$

(f) $p_2(n) = \lfloor n/2 \rfloor$ (floor function).

(f) $p_2(n) = \lfloor n/2 \rfloor$ is not negative.

Proof.

(a) A sum of positive integers is not negative.
(b) A partition into k parts has sum $\geq \underbrace{1+1+\dots+1}_{k \text{ times}} = k.$

(c) The only partition into 0 parts is $()$.

(d) The only partition of n into 1 part is (n) ,

which exists only if $n \neq 0,$

(e) Classify the partitions of n into k parts into 2 types:

Type 1: partitions whose last entry is 1.

Type 2: partitions whose last entry is $\neq 1$
(i.e., partitions whose all entries are ≥ 2).

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Now, there is a bijection

{Type-1 partitions of n into k parts}

\rightarrow {partitions of $n-1$ into $k-1$ parts},

$$(\lambda_1, \lambda_2, \dots, \lambda_{k-1}, 1) \mapsto (\lambda_1, \lambda_2, \dots, \lambda_{k-1}).$$

$$\text{Thus, } \#(\text{Type-1 partitions}) = p_{k-1}(n-1).$$

Also, there is a bijection

{Type-2 partitions of n into k parts}

\rightarrow {partitions of $n-k$ into k parts},

$$(\lambda_1, \lambda_2, \dots, \lambda_k) \mapsto (\lambda_1 - 1, \lambda_2 - 1, \dots, \lambda_k - 1).$$

$$\text{Thus, } \#(\text{Type-2 partitions}) = p_k(n-k).$$

Adding these equalities together, we get the claim of (e).

(f) The partitions of n into 2 parts are

$$(n-1, 1), (n-2, 2), (n-3, 3), \dots, \left(\left\lceil \frac{n}{2} \right\rceil, \left\lfloor \frac{n}{2} \right\rfloor\right).$$

↑
ceiling function

□

By the way ...

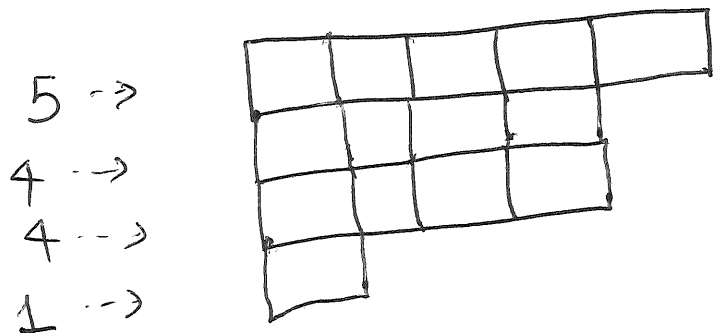
Prop. 4.13. Let $n \in \mathbb{N}$ and $k \in \mathbb{N}$. Then,

$p_k(n) =$ # of partitions of n whose largest part is k .

Proof outline. Picture proof: e.g., let $k=4$ and $n=14$.


Start with the partition $(5, 4, 4, 1)$ of n into k parts.

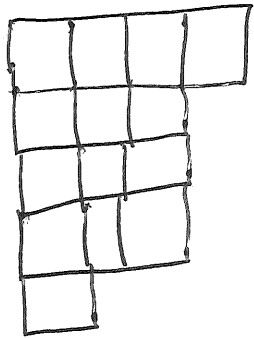
~~Draw~~ Draw a table of k left-aligned rows, where the length of each row is the corresponding part of the partition.



(This ~~table~~ table is called a Young

diagram or Ferrers diagram.)

Now, flip this table around the  diagonal, -9-



The length of the rows of this new table ~~are the~~ form
 a partition of n whose largest part is k .

This is a bijection.
 (It is called conjugation.)

(here, $(4, 3, 3, 3, 1)$.)

□

Prop. 4.14, (# of \uparrow $u \rightarrow u$ placements) = $\phi_{|X|}(|A|)$.
 surjective

Proof outline. Encode a surjective $u \rightarrow u$ placement as a partition
 of ~~the~~ $|A|$ into $|X|$ parts (where each part is the #
 of balls in some box). This is a bijection. □

Prop. 4.15. (# of $U \rightarrow U$ placements)

$$= p_0(|A|) + p_1(|A|) + \dots + p_{|X|}(|A|).$$

□

Proof. Exercise.

Prop. 4.16. (# of injective $U \rightarrow U$ placements) = $[|A| \leq |X|]$.

□

Proof. Exercise.

The twelvefold way: ~~twelvefold~~

f is...	arbitrary	injective	surjective
$A \rightarrow X$			
$L \rightarrow L$	$ X ^{ A }$	$ X (X -1)(X -2) \dots (X - A +1)$	$\text{sur}(A , X) \leftarrow = X ! \begin{Bmatrix} A \\ X \end{Bmatrix}$
$U \rightarrow L$	$\binom{ A + X -1}{ A }$	$\binom{ X }{ A }$	$\binom{ A -1}{ A - X }$
$L \rightarrow U$	$\begin{Bmatrix} A \\ 0 \end{Bmatrix} + \begin{Bmatrix} A \\ 1 \end{Bmatrix} + \dots + \begin{Bmatrix} A \\ X \end{Bmatrix}$	$[A \leq X]$	$\begin{Bmatrix} A \\ X \end{Bmatrix}$
$U \rightarrow U$	$p_0(A) + p_2(A) + \dots + p_{ X }(A)$	$[A \leq X]$	$p_{ X }(A)$

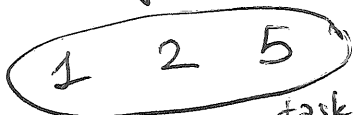
4.7, ODDS & ENDS

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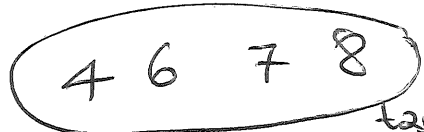
Example: Given n persons ($n > 0$) and k tasks ($k > 0$).

- (a) What is the # of ways to assign to each person a task such that each task has at least one person working on it?
- (b) What if we additionally want to choose a leader for each task (among the people assigned to it)?
- (c) What if, instead, we want to choose a vertical hierarchy (between the people working on the task) for each task?

Example: 8 people (1, 2, 3, 4, ..., 8), and 3 tasks.

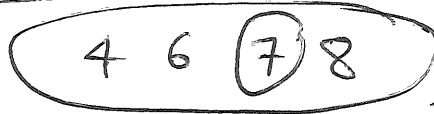
(a)  task 1

 task 2

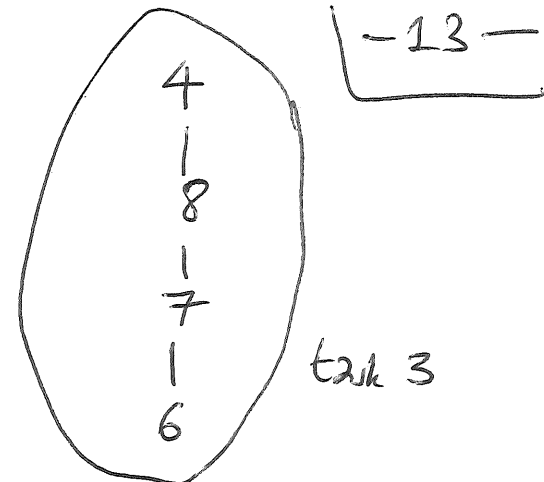
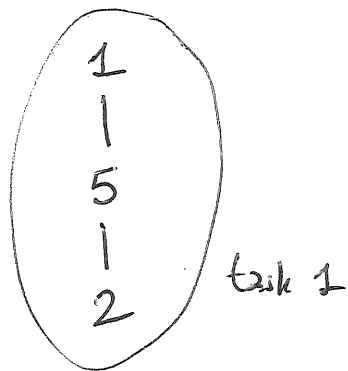
 task 3

(b)  task 1

 task 2

 task 3

(c)



Answers: (a) $\text{sur}(n, k)$,

(b) $(n(n-1)\cdots(n-k+1)) \cdot k^{n-k}$.

(c) $n! \cdot \binom{n-1}{k-1}$.

(see 4909 Fall 2017 Oct 10,)

4.8. A GLIMPSE OF RANKING & UNRANKING

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Ranking & unranking is an important subject; we only will scratch the surface. See [Loehr, Ch. 5] and [Knuth, "The Art of Computer Programming", vol. 4A] for more.

Statement of the problem: You have a finite set X . Find ways to

- (1) construct a list of all elements of X ("listing").
- (2) for each $k \in [1, |X|]$, compute the k -th element of this list without having to write down the whole list ("unranking").
- (3) for each $x \in X$, compute the rank of x in this list (i.e., the # of elements before x in the list) ("ranking").

Equivalently: construct a bijection $[1, |X|] \rightarrow X$ and its inverse. ("construct" = find an algorithm.)

Some questions apply to a countable set X , using $\{1, 2, 3, \dots\}$ instead $[1, |X|]$.

Example 1: let $N \in \mathbb{N}$. Let ~~$\mathcal{P}(\mathbb{N})$~~ $\mathcal{P}(X)$ (where X is any set) denote the set of all subsets of X .
How to list $\mathcal{P}(\mathbb{N})$ (a $2^{\mathbb{N}}$ -element set)?

Example: A list for $\mathcal{P}(\mathbb{N})$:
($\{1\}$, $\{2\}$, $\{3\}$, $\{1,2\}$, $\{2,3\}$, $\{1,3\}$, \emptyset , $\{1,2,3\}$).

This is not very systematic.

Better: recursion.

A list of $\mathcal{P}(\mathbb{N})$: $\{\emptyset\}$.
To build a list of $\mathcal{P}(\mathbb{N}+1)$ from a list of $\mathcal{P}(\mathbb{N})$,
we first write down this list of $\mathcal{P}(\mathbb{N})$,
and then write down a second copy of this list,
but replacing each subset S by $S \cup \{N+1\}$.

For example:

$\mathcal{P}(\mathbb{N})$	$\{\emptyset\}$
$\mathcal{P}(\mathbb{N}+1)$	$\{\emptyset, \{1\}\}$
$\mathcal{P}(\mathbb{N}+2)$	$\{\emptyset, \{1\}, \{2\}, \{1,2\}\}$
$\mathcal{P}(\mathbb{N}+3)$	$\{\emptyset, \{1\}, \{2\}, \{1,2\}, \{3\}, \{1,3\}, \{2,3\}, \{1,2,3\}\}$

Now, how do we rank/unrank in this list?

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Observations:

- The list of $\mathcal{P}([N+1])$ starts with the list of $\mathcal{P}([N])$,
Thus, the rank of a $S \in \mathcal{P}([N])$ does not depend on N ,
- If $S \in \mathcal{P}([N])$ has rank r ,
then $S \cup \{N+1\} \in \mathcal{P}([N+1])$ has rank $r + 2^N$,

\Rightarrow (by induction): The rank of any $S \in \mathcal{P}([N])$ is
$$\sum_{i \in S} 2^{i-1}$$

(E.g., the rank of $\{1,3\}$ is $2^{1-1} + 2^{3-1} = 1 + 4 = 5$.)

To unrank, we start with an $n \in \{0, 1, \dots, 2^N - 1\}$, and
we look for the subset S of $[N]$ with $\sum_{i \in S} 2^{i-1} = n$.
This is just the set of positions of the 1-bits in the
binary expansion of n .

Exercise: Ranking/unranking for lexicor subsets leads to
the Zeckendorf theorem.

Example 2: Given $j \in \mathbb{N}$ and ~~a~~ a set X , we let

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$\mathcal{P}_j(X)$ be the set of all j -element subsets of X ,
Given $j \in \mathbb{N}$, how to list $\mathcal{P}_j([N])$ or $\mathcal{P}_j(N)$?

Example: $j=3$; try to list $\mathcal{P}_j(\mathbb{N})$.

Bad attempt: ~~{0,1,2}~~, {0,1,3}, {0,1,4}, ...

This never gets to {0,2,3}.

Better: recursion. let's first list $\mathcal{P}_j(\{0,1,\dots,N\})$ for each ~~$N \geq -1$~~
 $N \geq -1$.

Again, $\mathcal{P}_0(\text{anything})$ has ^{list} (\emptyset) .

Also, $\mathcal{P}_j(\{0,1,\dots,-1\})$ has list $(\)$ if $j > 0$,

To build a list of $\mathcal{P}_j(\{0,1,\dots,N+1\})$ from a list of
 $\mathcal{P}_j(\{0,1,\dots,N\})$ and a list of $\mathcal{P}_{j-1}(\{0,1,\dots,N\})$, we
first write down the list of $\mathcal{P}_j(\{0,1,\dots,N\})$,
and then write down the list of $\mathcal{P}_{j-1}(\{0,1,\dots,N\})$,
but replacing each subset S by $S \cup \{N+1\}$.

For example:

$P_0(\emptyset)$	(\emptyset)
$P_0(\{0\})$	(\emptyset)
$P_1(\{0\})$	$(\{0\})$
$P_0(\{1\})$	(\emptyset)
$P_1(\{1\})$	$(\{0\}, \{1\})$
$P_2(\{1\})$	$(\{0, 1\})$
$P_0(\{0, 1, 2\})$	(\emptyset)
$P_1(\{0, 1, 2\})$	$(\{0\}, \{1\}, \{2\})$
$P_2(\{0, 1, 2\})$	$(\{0, 1\}, \{0, 2\}, \{1, 2\})$
$P_3(\{0, 1, 2\})$	$(\{0, 1, 2\})$
$P_4(\{0, 1, 2, 3\})$	$(\{0, 1\}, \{0, 2\}, \{0, 3\}, \{1, 2\}, \{1, 3\}, \{2, 3\})$

Again, the list of $\mathcal{P}_j(\{0, 1, \dots, N+1\})$ starts with the list of $\mathcal{P}_j(\{0, 1, \dots, N\})$. Thus, the rank of an $S \in \mathcal{P}_j(\{0, 1, \dots, N\})$ does not depend on N .
 How to rank/unrank?

Observe:

- If $S \in \mathcal{P}_{j-1}(\{0, 1, \dots, N\})$ has rank r , then $S \cup \{N+1\}$ has rank $\binom{N+1}{j} + r$.

\Rightarrow (by induction): The rank of any $\{s_1 < s_2 < \dots < s_j\} \in \mathcal{P}_j(\{0, 1, \dots, N\})$ is $\binom{s_1}{1} + \binom{s_2}{2} + \dots + \binom{s_j}{j}$.

This solves Exercise 7 on Midterm 1 again, as the list of $\mathcal{P}_j(\mathbb{N})$ is infinite in length when $j > 0$.

To unrank, you need to expand an $n \in \mathbb{N}$ into the form $n = \binom{s_1}{1} + \binom{s_2}{2} + \dots + \binom{s_j}{j}$ with $0 \leq s_1 < s_2 < \dots < s_j$.