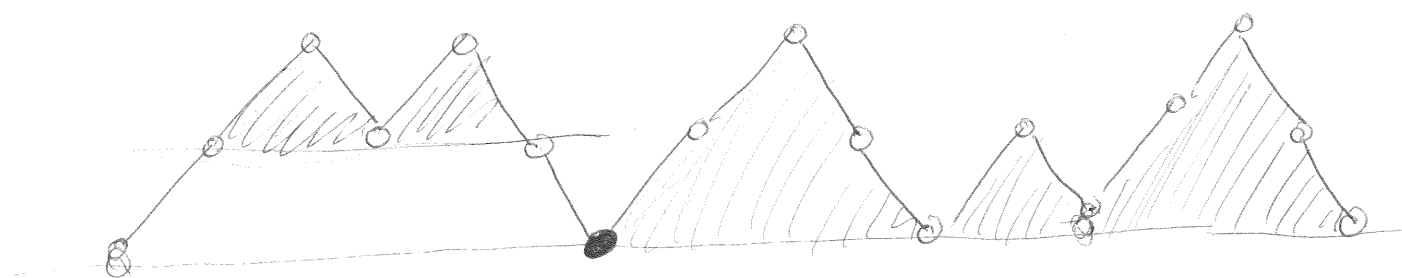


What is a recursion for c_n ?

(1)



Given $n > 0$, we want to see how many Dyck paths $(0,0) \rightarrow (2n,0)$ first return to the x -axis at $(2k,0)$ for a given $k \in [n]$.
(Note: The first return must be of the form $(2k,0)$ for some $k \in [n]$.)

To construct such a path, we

- first choose its first $2k$ steps:

they have to form a Dyck path that never returns to x -axis until $(2k,0)$;

in other words, they begin with a $\nearrow$, end with a $\searrow$, and inbetween proceed as a Dyck path ~~from~~ $(0,0) \rightarrow (2k-2,0)$.

$\Rightarrow c_{k-1}$ choices

(2)

• then choose its last $2(n-k)$ steps;

they have to proceed as a Dyck path
 $(0,0) \rightarrow (2(n-k), 0)$

$\Rightarrow c_{n-k}$ choices,

$$\Rightarrow c_n = \sum_{k=1}^n c_{k-1} c_{n-k}, \quad (1)$$

Let $C(x) = \sum_{n \geq 0} c_n x^n = c_0 + c_1 x + c_2 x^2 + \dots$

$$\begin{aligned} \text{Thus, } C(x) &= 1 + (c_0 c_0) x \\ &\quad + (c_0 c_1 + c_1 c_0) x^2 \\ &\quad + (c_0 c_2 + c_1 c_1 + c_2 c_0) x^3 \\ &\quad + \dots \\ &= 1 + x \underbrace{(c_0 + c_1 x + c_2 x^2 + \dots)}_{=C(x)}^2, \end{aligned}$$

so $C(x) = 1 + x (C(x))^2$.

This is a quadratic in $C(x)$. Solve it (whenever it's allowed). Get ④

$$C(x) = \frac{1 \pm \sqrt{1-4x}}{2x}.$$

The $\pm$ cannot be $+$, since with $+$ the power series on top of the fraction would not be divisible by $2x$. So

$$C(x) = \frac{1 - \sqrt{1-4x}}{2x} = \cancel{\frac{1}{2x}} \frac{1}{2x} (1 - (1-4x)^{1/2}).$$

But recall $(1+x)^n = \sum_{k=0}^n \binom{n}{k} x^k = \sum_k \binom{n}{k} x^k.$

Let's pretend this works for $n = 1/2$, too.

$$\text{Thus, } (1 + \cancel{x})^{1/2} = \sum_k \binom{1/2}{k} x^k.$$

Now, substitute $-4x$ for x here, to get

$$\begin{aligned} (1-4x)^{1/2} &= \sum_k \binom{1/2}{k} \cancel{(-4x)}^k \\ &= \sum_k \binom{1/2}{k} (-4)^k x^k. \end{aligned}$$

So our formula for $C(x)$ becomes

(5)

$$C(x) = \frac{1}{2x} \left(1 - \sum_k \binom{1/2}{k} (-4)^k x^k \right)$$

$$= - \sum_{k \geq 1} \binom{1/2}{k} (-4)^k x^k$$

$$= - \sum_{k \geq 1} \binom{1/2}{k} \frac{(-4)^k}{2} x^{k-1}$$

$$= - \sum_{k \geq 0} \binom{1/2}{k+1} \frac{(-4)^{k+1}}{2} x^k$$

Now, comparing coeffs before x^n gives

$$C_n = - \binom{1/2}{n+1} \frac{(-4)^{n+1}}{2}$$

$$= \frac{1/2}{n+1} \binom{-1/2}{n}$$

$$= \binom{2n}{n} \left(-\frac{1}{4}\right)^n \quad (\text{by HW2 ex 1})$$

$$= \cancel{\frac{1/2}{n+1}} \binom{2n}{n} \cancel{\left(-\frac{1}{4}\right)^n} \frac{\cancel{(-4)^{n+1}}}{\cancel{2}}$$

$$= \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n-1}$$

Ex. 3, Thm. 3.12 (= Thm 12 in ch. 3) says (6)
 that $\binom{a+b}{n} = \sum_{k=0}^n \binom{a}{k} \binom{b}{n-k} \quad \forall a, b \in \mathbb{N},$
 $n \in \mathbb{N}.$

We can reproove it using FPS.

$$\begin{aligned}
 \text{Namely, } (1+x)^a (1+x)^b &= (1+x)^{a+b} \\
 &= \underbrace{\sum_k \binom{a}{k} x^k}_{= \sum_k \binom{a}{k} x^k} \underbrace{\sum_l \binom{b}{l} x^l}_{= \sum_l \binom{b}{l} x^l} = \sum_k \binom{a+b}{k} x^k \\
 &= \left(\sum_k \binom{a}{k} x^k \right) \left(\sum_l \binom{b}{l} x^l \right) \\
 &= \sum_{k,l} \binom{a}{k} x^k \binom{b}{l} x^l \\
 &= \sum_{k+l} \binom{a}{k} \binom{b}{l} x^{k+l} \\
 &= \sum_n \left(\sum_{k=0}^n \binom{a}{k} \binom{b}{n-k} \right) x^n
 \end{aligned}$$

$$\text{So } \sum_n \left(\sum_{k=0}^n \binom{a}{k} \binom{b}{n-k} \right) x^n = \sum_n \binom{a+b}{n} x^n.$$

$$\text{Comparing coeffs yields } \sum_{k=0}^n \binom{a}{k} \binom{b}{n-k} = \binom{a+b}{n}.$$

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Ex. 4. (from Wilf's "generatingfunctionology")

Define a sequence $(a_0, a_1, a_2, \dots)$ recursively by

$$a_0 = 1, \quad a_{n+1} = 2a_n + n \quad \forall n \geq 0.$$

(So it starts out 1, 2, 5, 12, 27, 58, 121, ...)

(OEIS: A000325, shifted.)

$$\text{Set } A(x) = a_0 + a_1 x + a_2 x^2 + \dots$$

$$\text{Now, } A(x) = 1 + (2a_0 + 0)x + (2a_1 + 1)x^2 + (2a_2 + 2)x^3 + \dots$$

$$= 1 + 2(a_0 + a_1 x + a_2 x^2 + \dots)$$

$$+ (0x + 1x^2 + 2x^3 + \dots)$$

$$= 1 + 2x A(x) + x(0 + 1x + 2x^2 + 3x^3 + \dots) = ?$$

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Now, $0 + 1x + 2x^2 + 3x^3 + \dots = ?$

Two ways to compute it:

$$\begin{aligned}
 \textcircled{1} \quad & 0 + 1x + 2x^2 + 3x^3 + \dots = x(1 + 2x + 3x^2 + \dots) \\
 & \qquad \qquad \qquad = \underbrace{\left(\cancel{x} 1 + x + x^2 + \dots \right)'} \\
 & \qquad \qquad \qquad = \left(\frac{1}{1-x} \right)' = \frac{1}{(1-x)^2}
 \end{aligned}$$

$$\begin{aligned}
 & = x \cdot \frac{1}{(1-x)^2} \\
 \textcircled{2} \quad & 0 + 1x + 2x^2 + 3x^3 + \dots = x + x^2 + x^3 + x^4 + \dots \\
 & \qquad \qquad \qquad + x^2 + x^3 + x^4 + \dots \\
 & \qquad \qquad \qquad + x^3 + x^4 + \dots \\
 & \qquad \qquad \qquad + x^4 + \dots \\
 & \qquad \qquad \qquad + \dots
 \end{aligned}$$

$$= x \cdot \frac{1}{1-x} + x^2 \cdot \frac{1}{1-x} + x^3 \cdot \frac{1}{1-x} + \dots$$

$$= \underbrace{\left(x + x^2 + x^3 + \dots\right)}_{= x \cdot \frac{1}{1-x}} \cdot \frac{1}{1-x}$$

$$= x \cdot \left(\frac{1}{1-x}\right)^2$$

[But be careful with ∞ sums.]

So our formula for $A(x)$ becomes

$$A(x) = 1 + 2x A(x) + x \cdot \frac{1}{(1-x)^2}$$

This is a linear eqn. in $A(x)$. Solving it yields

$$A(x) = \frac{1 - 2x + 2x^2}{(1-x)^2(1-2x)} \stackrel{\text{PFD}}{=} \underbrace{\frac{-1}{(1-x)^2}} + \underbrace{\frac{2}{1-2x}} = -\underbrace{(1+2x+3x^2+\dots)}_{\text{(by above)}} \bigg| = 2 \sum_k (2x)^k = \sum_k 2^{k+1} x^k$$

$$= -(1+2x+3x^2+\dots) + \sum_k 2^{k+1} x^k = \sum_k (2^{k+1} - (k+1)) x^k$$

Comparing coefficients yields $a_n = 2^{n+1} - (n+1)$. (10)

4.2. THE CONCEPT OF A FPS (FORMAL POWER SERIES)

We want to give a definition of a FPS and also justifications for the things we did to FPS (dividing, quadratic eqns, ∞ sums, ...).

First things first: FPS are not functions. You cannot substitute $x=2$ into $\frac{1}{1-x} = 1+x+x^2+\dots$ and "obtain" $\frac{1}{-1} = 1+2+4+8+16+\dots$.

Let us go back to abstract algebra.

Def. A commutative ring $\mathbb{B}$, informally, a set $\mathbb{B}$ equipped with

binary operations $\oplus, \ominus$, and $\odot$ and elements $\odot$ and $\odot$ that "behave" like $\mathbb{N}$ addition, subtraction, multiplication of numbers, and the numbers 0 and 1, respectively.

For example, they have to satisfy rules like ~~look~~ _{mean}

$$(a \oplus b) \odot c = a \odot c \oplus b \odot c.$$

Formally, a comm. ring is a set K equipped with

maps $\oplus: K \times K \rightarrow K$, $\ominus: K \times K \rightarrow K$, $\odot: K \times K \rightarrow K$

and elts $\odot \in K$ and $\mathbb{1} \in K$ satisfying the following axioms.

(a) Commutativity of $\oplus$: $a \oplus b = b \oplus a$.

(b) Associativity of $\oplus$: $a \oplus (b \oplus c) = (a \oplus b) \oplus c$.

(c) Neutrality of $\odot$: $a \oplus \odot = \odot \oplus a = a$.

(d) $\ominus$ undoes $\oplus$: $a \oplus b = c \iff a = c \ominus b$.

(e) Commutativity of $\ominus$: $a \ominus b = b \ominus a$.

(f) Associativity of $\ominus$: $a \ominus (b \ominus c) = (a \ominus b) \ominus c$.

(g) Distributivity: $a \odot (b \oplus c) = (a \odot b) \oplus (a \odot c)$;
 $(a \oplus b) \odot c = (a \odot c) \oplus (b \odot c)$

(h) Neutrality of $\mathbb{1}$: $a \odot \mathbb{1} = \mathbb{1} \odot a = a$.

(i) Annihilation: $a \odot \odot = \odot \odot a = \odot$.

Examples: $\mathbb{Z}, \mathbb{Q}, \mathbb{R}, \mathbb{C}$.

Non-examples: N (semiring); $\mathbb{Q}^{m \times m}$ (noncomm. ring), for $m > 1$.

Most authors do this part differently, but the end-result is the same.

Usually, drop the circles, so write $\oplus, \ominus, \otimes, \oslash, \pm, \mp$.

(12)

See [detnotes], start of ch. 5.

Good news: In an arbitrary commutative ring, the standard

rules of computation apply:

— You can compute finite sums without caring about the order of summation or parenthesization!

$$((a + (b + c)) + d) + e = (a + b) + (c + (d + e))$$

("general associativity"),

so you ~~can~~ just write $a + b + c + d + e$,

Also, $a + b + c + d + e = d + b + a + e + c$

("general commutativity").

Ex 0. Counting compositions using FPS.

①

Let $k(n)$ be the # of compositions of n .
Want $k(n)$.

$$\text{Set } K(x) = k(0) + k(1)x + k(2)x^2 + \dots$$

$$\text{Thus, } K(x) = \sum_{n \geq 0} k(n)x^n = \sum_{n \geq 0} \sum_{\substack{(a_1, \dots, a_k) \\ \text{comp of } n}} x^n,$$

$$= \sum_{\substack{(a_1, a_2, \dots, a_k) \\ \text{composition} \\ \text{(i.e. tuple of} \\ \text{pos ints)}}} x^{a_1 + a_2 + \dots + a_k}$$

$$= \sum_{k \geq 0} \sum_{(a_1, \dots, a_k) \in \{1, 2, 3, \dots\}^k} x^{a_1 + a_2 + \dots + a_k}$$

$$\stackrel{\text{product rule}}{=} \left(\sum_{a=1}^{\infty} x^a \right)^k$$

$$= \sum_{k \geq 0} \left(\sum_{a=1}^{\infty} x^a \right)^k = \frac{1}{1 - \sum_{a=1}^{\infty} x^a}$$

$$\approx \quad \left(\text{since } \sum_{k \geq 0} y^k = \frac{1}{1-y} \right)$$

$$= \frac{1}{1 - \left(\frac{1}{1-x} - 1 \right)} \quad \left(\text{since } \sum_{a=1}^{\infty} x^a = \sum_{a \geq 0} x^a - 1 = \frac{1}{1-x} - 1 \right)$$

$$= \frac{1}{2 - \frac{1}{1-x}} = \frac{1-x}{2(1-x)-1} = \frac{1-x}{1-2x}$$

$$= \frac{1}{1-2x} - x \cdot \frac{1}{1-2x} = 1 + \frac{x}{1-2x}$$

$$= 1 + x \sum_{k \geq 0} (2x)^k \quad \left(\text{since } \frac{1}{1-y} = \sum_{k \geq 0} y^k \right)$$

$$= 1 + \sum_k 2^k x^{k+1} = \sum_k \begin{cases} 2^{k-1} & \text{if } k \geq 1 \\ 1 & \text{if } k=0 \end{cases} \cdot x^k$$

Compare coeffs

$$\Rightarrow k(n) = \begin{cases} 2^{n-1} & \text{if } n > 0 \\ 1 & \text{if } n = 0 \end{cases}$$