## An exercise on source and sink mutations of acyclic quivers\*

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In this note, we will use the following notations (which come from Lampe's notes [Lampe, §2.1.1]):

• A *quiver* means a tuple  $Q = (Q_0, Q_1, s, t)$ , where  $Q_0$  and  $Q_1$  are two finite sets and where s and t are two maps from  $Q_1$  to  $Q_0$ . We call the elements of  $Q_0$  the *vertices* of the quiver Q, and we call the elements of  $Q_1$  the *arrows* of the quiver Q. For every  $e \in Q_1$ , we call s(e) the *starting point* of e (and we say that e starts at s(e)), and we call t(e) the *terminal point* of e (and we say that e ends at t(e)). Furthermore, if  $e \in Q_1$ , then we say that e is an arrow from s(e) to t(e).

So the notion of a quiver is one of many different versions of the notion of a finite directed graph. (Notice that it is a version which allows multiple arrows, and which distinguishes between them – i.e., the quiver stores not just the information of how many arrows there are from a vertex to another, but it actually has them all as distinguishable objects in  $Q_1$ . Lampe himself seems to later tacitly switch to a different notion of quivers, where edges from a given to vertex to another are indistinguishable and only exist as a number. This does not matter for the next exercise, which works just as well with either notion of a quiver; but I just wanted to have it mentioned.)

• The underlying undirected graph of a quiver  $Q = (Q_0, Q_1, s, t)$  is defined as the undirected multigraph with vertex set  $Q_0$  and edge multiset

$$\{\{s(e),t(e)\}\mid e\in Q_1\}_{\text{multiset}}.$$

("Multigraph" means that multiple edges are allowed, but we do not make them distinguishable.)

<sup>\*</sup>This used to be Chapter 7 of my notes "Notes on the combinatorial fundamentals of algebra" (version of 7 November 2018), but has since been removed from the latter notes.

- A quiver  $Q = (Q_0, Q_1, s, t)$  is said to be *acyclic* if there is no sequence  $(a_0, a_1, \ldots, a_n)$  of elements of  $Q_0$  such that n > 0 and  $a_0 = a_n$  and such that Q has an arrow from  $a_i$  to  $a_{i+1}$  for every  $i \in \{0, 1, \ldots, n-1\}$ . (This is equivalent to [Lampe, Definition 2.1.7].) Notice that this does not mean that the *underlying undirected graph* of Q has no cycles.
- Let  $Q = (Q_0, Q_1, s, t)$ . Then, a *sink* of Q means a vertex  $v \in Q_0$  such that no  $e \in Q_1$  starts at v (in other words, no arrow of Q starts at v). A *source* of Q means a vertex  $v \in Q_0$  such that no  $e \in Q_1$  ends at v (in other words, no arrow of Q ends at v).
- Let  $Q = (Q_0, Q_1, s, t)$ . If  $i \in Q_0$  is a sink of Q, then the *mutation*  $\mu_i(Q)$  of Q at i is the quiver obtained from Q simply by turning all arrows ending at i. (To be really pedantic: We define  $\mu_i(Q)$  as the quiver  $(Q_0, Q_1, s', t')$ , where

$$s'\left(e\right) = \begin{cases} t\left(e\right), & \text{if } t\left(e\right) = i; \\ s\left(e\right), & \text{if } t\left(e\right) \neq i \end{cases} \qquad \text{for each } e \in Q_1$$
 and 
$$t'\left(e\right) = \begin{cases} s\left(e\right), & \text{if } t\left(e\right) = i; \\ t\left(e\right), & \text{if } t\left(e\right) \neq i \end{cases} \qquad \text{for each } e \in Q_1.$$

) If  $i \in Q_0$  is a source of Q, then the *mutation*  $\mu_i(Q)$  of Q at i is the quiver obtained from Q by turning all arrows starting at i. (Notice that if i is both a source and a sink of Q, then these two definitions give the same result; namely,  $\mu_i(Q) = Q$  in this case.)

If Q is an acyclic quiver, then  $\mu_i(Q)$  is acyclic as well (whenever  $i \in Q_0$  is a sink or a source of Q).

We use the word "mutation" not only for the quiver  $\mu_i(Q)$ , but also for the operation that transforms Q into  $\mu_i(Q)$ . (We have defined this operation only if i is a sink or a source of Q. It can be viewed as a particular case of the more general definition of mutation given in [Lampe, Definition 2.2.1], at least if one gives up the ability to distinguish different arrows from one vertex to another.)

<sup>&</sup>lt;sup>1</sup>To *turn* an arrow e means to reverse its direction, i.e., to switch the values of s(e) and t(e). We model this as a change to the functions s and t, not as a change to the arrow itself.

**Exercise 0.1.** Let  $Q = (Q_0, Q_1, s, t)$  be an acyclic quiver.

(a) Let A and B be two subsets of  $Q_0$  such that  $A \cap B = \emptyset$  and  $A \cup B = Q_0$ . Assume that there exists no arrow of Q that starts at a vertex in B and ends at a vertex in A. Then, by turning all arrows of Q which start at a vertex in A and end at a vertex in B, we obtain a new acyclic quiver  $\text{mut}_{A,B} Q$ .

(When we say "turning all arrows of Q which start at a vertex in A and end at a vertex in B", we mean "turning all arrows e of Q which satisfy  $s(e) \in A$  and  $t(e) \in B$ ". We do **not** mean that we fix a vertex e in A and a vertex e in B, and only turn the arrows from e to e.)

For example, if 
$$Q = 3 \longrightarrow 4$$
 and  $A = \{1,3\}$  and  $B = \{2,4\}$ , then

Prove that  $\operatorname{mut}_{A,B} Q$  can be obtained from Q by a sequence of mutations at sinks. (More precisely, there exists a sequence  $\left(Q^{(0)},Q^{(1)},\ldots,Q^{(\ell)}\right)$  of acyclic quivers such that  $Q^{(0)}=Q$ ,  $Q^{(\ell)}=\operatorname{mut}_{A,B}Q$ , and for every  $i\in\{1,2,\ldots,\ell\}$ , the quiver  $Q^{(i)}$  is obtained from  $Q^{(i-1)}$  by mutation at a sink of  $Q^{(i-1)}$ .)

[In our above example, we can mutate at 4 first and then at 2.]

- **(b)** If  $i \in Q_0$  is a **source** of Q, then show that the mutation  $\mu_i(Q)$  can be obtained from Q by a sequence of mutations at sinks.
- (c) Assume now that the underlying **undirected** graph of Q is a tree. (In particular, Q cannot have more than one edge between two vertices, as these would form a cycle in the underlying undirected graph!) Show that any acyclic quiver which can be obtained from Q by turning some of its arrows can also be obtained from Q by a sequence of mutations at sinks.

**Remark 0.1.** More general results than those of Exercise 0.1 are stated (for directed graphs rather than quivers, but it is easy to translate from one language into another) in [Pretzel]. An equivalent version of Exercise 0.1 **(c)** also appears as Exercise 6 in [GrRaOg] (because a quiver *Q* whose underlying undirected graph is a tree can be regarded as an orientation of the latter tree, and because the concept of "pushing sources" in [GrRaOg] corresponds precisely to our concept of mutations at sinks, except that all arrows need to be reversed).

Solution to Exercise 0.1. (a) We prove the claim by induction over |B|.

*Induction base:* Assume that |B| = 0. Thus,  $B = \emptyset$ , and thus there are no arrows of Q which start at a vertex in A and at a vertex in B. Hence,  $\operatorname{mut}_{A,B} Q = Q$ , and this can clearly be obtained from Q by a sequence of mutations at sinks (namely, by the empty sequence). Thus, Exercise 0.1 (a) holds if |B| = 0. This completes the

induction base.<sup>2</sup>

*Induction step:*<sup>3</sup> Let  $N \in \mathbb{N}$ . Assume that Exercise 0.1 (a) holds whenever |B| = N. We now need to prove that Exercise 0.1 (a) holds whenever |B| = N + 1.

So let A and B be two subsets of  $Q_0$  such that  $A \cap B = \emptyset$  and  $A \cup B = Q_0$ . Assume that there exists no arrow of Q that starts at a vertex in B and ends at a vertex in A. Assume further that |B| = N + 1. We need to prove that  $\operatorname{mut}_{A,B} Q$  can be obtained from Q by a sequence of mutations at sinks.

Notice that  $B = Q_0 \setminus A$  (since  $A \cap B = \emptyset$  and  $A \cup B = Q_0$ ).

It is easy to see that there exists some  $b \in B$  such that

there is no 
$$e \in Q_1$$
 satisfying  $t(e) = b$  and  $s(e) \in B$  (1)

<sup>4</sup>. Fix such a *b*. Clearly,  $b \notin A$  (since  $b \in B = Q_0 \setminus A$ ).

Now,  $A \cup \{b\}$  and  $B \setminus \{b\}$  are two subsets of  $Q_0$  such that  $(A \cup \{b\}) \cap (B \setminus \{b\}) = \emptyset$  and  $(A \cup \{b\}) \cup (B \setminus \{b\}) = Q_0$  5. Furthermore, there exists no arrow of Q that starts at a vertex in  $B \setminus \{b\}$  and ends at a vertex in  $A \cup \{b\}$  6. Hence,  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$  is a well-defined acyclic quiver. Moreover, since  $b \in B$ , we have  $|B \setminus \{b\}| = \bigcup_{=N+1}^{B} (-1) = N + 1 - 1 = N$ . Thus, Exercise 0.1 (a) can be applied

to  $A \cup \{b\}$  and  $B \setminus \{b\}$  instead of A and B (by the induction hypothesis). As a consequence, we conclude that  $\min_{A \cup \{b\}, B \setminus \{b\}} Q$  can be obtained from Q by a sequence of mutations at sinks.

<sup>&</sup>lt;sup>2</sup>Yes, this was a completely honest induction base. You don't need to start at |B| = 1 unless you want to use something like |B| > 1 in the induction step (but even then, you should also handle the case |B| = 0 separately).

<sup>&</sup>lt;sup>3</sup>The letter  $\mathbb N$  denotes the set  $\{0,1,2,\ldots\}$  here.

<sup>&</sup>lt;sup>4</sup>*Proof.* Assume the contrary. Thus, for every  $b \in B$ , there is an  $e \in Q_1$  satisfying t(e) = b and  $s(e) \in B$ . Let us fix such an e (for each  $b \in B$ ), and denote it by  $e_b$ .

Thus, for every  $b \in B$ , we have  $e_b \in Q_1$  and  $t(e_b) = b$  and  $s(e_b) \in B$ . We can thus define a sequence  $(b_0, b_1, b_2, \ldots)$  of vertices in B recursively as follows: Set  $b_0 = b$ , and set  $b_{i+1} = s(e_{b_i})$  for every  $i \in \mathbb{N}$ . Thus,  $(b_0, b_1, b_2, \ldots)$  is an infinite sequence of elements of B. Since B is a finite set, this sequence must thus pass through an element twice (to say the least). In other words, there are two positive integers u and v such that u < v and  $b_u = b_v$ . Consider these u and v.

Now, for every  $i \in \mathbb{N}$ , we have  $t\left(e_{b_i}\right) = b_i$  (by the definition of  $e_{b_i}$ ) and  $s\left(e_{b_i}\right) = b_{i+1}$ . Thus, for every  $i \in \mathbb{N}$ , the arrow  $e_{b_i}$  is an arrow from  $b_{i+1}$  to  $b_i$ . Thus, there is an arrow from  $b_{i+1}$  to  $b_i$  for every  $i \in \mathbb{N}$ . In particular, we have an arrow from  $b_v$  to  $b_{v-1}$ , an arrow from  $b_{v-1}$  to  $b_{v-2}$ , etc., and an arrow from  $b_{u+1}$  to  $b_u$ . Since  $b_u = b_v$ , these arrows form a cycle in Q, which contradicts the hypothesis that the quiver Q is acyclic. This contradiction proves that our assumption was wrong, qed.

<sup>&</sup>lt;sup>5</sup>*Proof.* These are easy exercises in set theory. Use  $A \cap B = \emptyset$  and  $A \cup B = Q_0$  and  $b \in B$ .

<sup>&</sup>lt;sup>6</sup>*Proof.* Assume the contrary. Thus, there exists an arrow of Q that starts at a vertex in  $B \setminus \{b\}$  and ends at a vertex in  $A \cup \{b\}$ . Let e be such an arrow. Then,  $s(e) \in B \setminus \{b\}$  and  $t(e) \in A \cup \{b\}$ .

We have  $s(e) \in B \setminus \{b\} \subseteq B$ . Thus,  $t(e) \neq b$  (because having t(e) = b would contradict (1)). Combined with  $t(e) \in A \cup \{b\}$ , this yields  $t(e) \in (A \cup \{b\}) \setminus \{b\} \subseteq A$ . Thus, e is an arrow of Q that starts at a vertex in B (since  $s(e) \in B$ ) and ends at a vertex in A (since  $t(e) \in A$ ). This contradicts our hypothesis that there exists no arrow of Q that starts at a vertex in B and ends at a vertex in A. This is the desired contradiction, and so we are done.

We shall now prove that  $\operatorname{mut}_{A,B} Q$  can be obtained from  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$  by a mutation at a sink. In fact, b is a sink of  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$ . Hence, the mutation  $\mu_b \left( \operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q \right)$  is well-defined. We now have

$$\operatorname{mut}_{A,B} Q = \mu_b \left( \operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q \right) \tag{2}$$

<sup>8</sup>. Therefore,  $\operatorname{mut}_{A,B} Q$  can be obtained from  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$  by a single mutation

<sup>7</sup>*Proof.* Assume the contrary. Thus, there exists an arrow e of  $\text{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$  which starts at b. Consider this e.

Recall that  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  was obtained from Q by turning all arrows of Q which start at a vertex in  $A\cup\{b\}$  and end at a vertex in  $B\setminus\{b\}$ . Thus, every arrow of  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  which starts at a vertex in  $B\setminus\{b\}$  and ends at a vertex in  $A\cup\{b\}$  has originally been going in the opposite direction in Q (because there exists no arrow of Q that starts at a vertex in  $B\setminus\{b\}$  and ends at a vertex in  $A\cup\{b\}$ , while all the other arrows of  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  have been copied over unchanged from Q. The arrow e of  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  starts at e (which is not an element of e in e in

Recall that there exists no arrow of Q that starts at a vertex in B and ends at a vertex in A. Thus, an arrow of Q which starts at a vertex in B must not end at a vertex in A. In particular, the arrow e of Q must not end at a vertex in A (because it starts at  $b \in B$ ). Hence, the arrow e of Q ends at a vertex in  $Q_0 \setminus A = B$ . In other words,  $t(e) \in B$ .

We cannot have t(e) = s(e) (because otherwise, the arrow e would form a cycle, but the quiver Q is acyclic). Hence,  $t(e) \neq s(e) = b$  (since e starts at b). Combined with  $t(e) \in B$ , this yields  $t(e) \in B \setminus \{b\}$ .

Thus, the arrow e of Q starts at a vertex in  $A \cup \{b\}$  (since  $s(e) = b \in A \cup \{b\}$ ) and ends at a vertex in  $B \setminus \{b\}$  (since  $t(e) \in B \setminus \{b\}$ ). As we know,  $\max_{A \cup \{b\}, B \setminus \{b\}} Q$  was obtained from Q by turning all such arrows. Hence, the arrow e must have been turned when it became an arrow of  $\max_{A \cup \{b\}, B \setminus \{b\}} Q$ . But this contradicts the fact that the arrow e has been copied over unchanged from Q. This contradiction proves that our assumption was wrong, qed.

<sup>8</sup>Proof of (2): We have  $Q_0 = A \cup \underbrace{B}_{=\{b\} \cup (B \setminus \{b\})} = A \cup \{b\} \cup (B \setminus \{b\}).$ 

Recall that the quiver  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  was obtained from Q by turning all arrows of Q which start at a vertex in  $A\cup\{b\}$  and end at a vertex in  $B\setminus\{b\}$ . Furthermore, the quiver  $\mu_b\left(\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q\right)$  was obtained from  $\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q$  by turning all arrows ending at b. Thus,  $\mu_b\left(\operatorname{mut}_{A\cup\{b\},B\setminus\{b\}}Q\right)$  can be obtained from Q by a two-step process, where

- in the first step, we turn all arrows of Q which start at a vertex in  $A \cup \{b\}$  and end at a vertex in  $B \setminus \{b\}$ ;
- in the second step, we turn all arrows ending at *b*.

Now, let us analyze what this two-step process does to an arrow of *Q*, depending on where the arrow starts and ends:

1. If *e* is an arrow of *Q* which ends at a vertex in *A*, then this arrow never gets turned during our process. Indeed, let *e* be such an arrow. Then, *e* ends at a vertex in *A*, and thus does

at a sink (namely, at the sink b). Since  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$  (in turn) can be obtained from Q by a sequence of mutations at sinks, this shows that  $\operatorname{mut}_{A,B} Q$  can be obtained from Q by a sequence of mutations at sinks (namely, we first need to mutate

not end at a vertex in B (since  $A \cap B = \emptyset$ ); therefore, it does not end at a vertex in  $B \setminus \{b\}$  either. Hence, the first step does not turn it. Therefore, after the first step, it still does not end at a vertex in B (since it did not end at a vertex in B originally). In particular, it does not end at b (since  $b \in B$ ). Hence, it does not get turned at the second step either. So, b never turns, and thus retains its original direction throughout the process.

- 2. If *e* is an arrow of *Q* which ends at *b*, then this arrow gets turned once (namely, at the second step). Thus, its direction is reversed at the end of the process.
- 3. If e is an arrow of Q which starts at a vertex in A and ends at a vertex in  $B \setminus \{b\}$ , then this arrow gets turned once (namely, at the first step). Here is why: Let e be an arrow of Q which starts at a vertex in A and ends at a vertex in  $B \setminus \{b\}$ . Then, e starts at a vertex in  $A \cup \{b\}$  and ends at a vertex in  $B \setminus \{b\}$  (since  $A \subseteq A \cup \{b\}$ ). Thus, it gets turned at the first step. After this, it becomes an arrow which ends at a vertex in A (because originally it started at a vertex in A), and so it does not end at b (because  $b \notin A$ ). Therefore, it does not turn at the second step; hence, it has turned exactly once altogether. Its direction is therefore reversed at the end of the process.
- 4. If e is an arrow of Q which starts at b and ends at a vertex in  $B \setminus \{b\}$ , then this arrow gets turned twice (once at each step). Indeed, let e be such an arrow. Then, e starts at a vertex in  $A \cup \{b\}$  (namely, at b) and ends at a vertex in  $B \setminus \{b\}$ . Hence, it gets turned at the first step. After that, it ends at b (because it used to start at b before it was turned), and therefore it gets turned again at the second step. Hence, the direction of e at the end of the two-step process is again the same as it was in Q.
- 5. If e is an arrow of Q which starts at a vertex in  $B \setminus \{b\}$  and ends at a vertex in  $B \setminus \{b\}$ , then this arrow never gets turned. Indeed, it starts at a vertex in  $B \setminus \{b\}$ ; thus, it does **not** start at a vertex in  $A \cup \{b\}$  (since  $B \setminus \{b\} = (Q_0 \setminus A) \setminus \{b\} = Q_0 \setminus (A \cup \{b\})$ ). Hence,

it does not get turned at the first step. Moreover, in Q, this arrow e does not end at b (because it ends at a vertex in  $B \setminus \{b\}$ ); thus it does not end at b after the first step either (since it does not get turned at the first step). Hence, it does not get turned at the second step either. Therefore, e never gets turned, and thus retains its original direction from Q after the two-step process.

The five cases we have just considered cover all possibilities (because every arrow e either ends at a vertex in A or ends at b or ends at a vertex in  $B \setminus \{b\}$ ; and in the latter case, it either starts at a vertex in A, or starts at b, or starts at a vertex in  $B \setminus \{b\}$  (since  $Q_0 = A \cup \{b\} \cup (B \setminus \{b\})$ )). From our case analysis, we can draw the following conclusions:

- If *e* is an arrow of *Q* which starts at a vertex in *A* and ends at a vertex in *B*, then the arrow *e* has reversed its orientation at the end of the two-step process. (This follows from our Cases 2 and 3 above.)
- If *e* is an arrow of *Q* which starts at a vertex in *B* or ends at a vertex in *A*, then this arrow *e* has the same orientation at the end of the two-step process as it did in *Q*. (Indeed, let us prove this. Let *e* be an arrow of *Q* which starts at a vertex in *B* or ends at a vertex in *A*. We need to show that *e* has the same orientation at the end of the two-step process as it did in *Q*. If *e* ends at a vertex in *A*, then this follows from our analysis of Case 1. So let us assume that *e* does not end at a vertex in *A*. Hence, *e* must start at a vertex in *B* (since *e* starts at a vertex in *B* or ends at a vertex in *A*). In other words, *s* (*e*) ∈ *B*. Hence,

at the sinks that give us  $\operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q$ , and then we have to mutate at b). This proves that Exercise 0.1 (a) holds whenever |B| = N + 1. The induction step is complete, and thus Exercise 0.1 (a) is solved.

**(b)** Let  $i \in Q_0$  be a source in Q. Let  $A = \{i\}$  and  $B = Q_0 \setminus A$ . Then, A and B are two subsets of  $Q_0$  such that  $A \cap B = \emptyset$  and  $A \cup B = Q_0$ . There exists no arrow of Q that starts at a vertex in B and ends at a vertex in  $A^{-9}$ . Hence, the quiver  $\operatorname{mut}_{A,B} Q$  is well-defined. Moreover, this quiver  $\operatorname{mut}_{A,B} Q$  is obtained by turning all arrows of Q which start at a vertex in A and end at a vertex in B. But these arrows are precisely the arrows of Q starting at  $i^{-10}$ . Hence,  $\operatorname{mut}_{A,B} Q$  is obtained by turning all arrows of Q starting at  $i^{-10}$ . Hence,  $\operatorname{mut}_{A,B} Q$  is obtained by turning all arrows of Q starting at Q. Now, Exercise 0.1 (a) shows that  $\operatorname{mut}_{A,B} Q$  can be obtained from Q by a sequence of mutations at sinks. Hence,  $\mu_i(Q)$  can be obtained from Q by a sequence of mutations at sinks (since  $\operatorname{mut}_{A,B} Q = \mu_i(Q)$ ). Exercise 0.1 (b) is proven.

(c) Let Q' be any acyclic quiver which can be obtained from Q by turning some of its arrows. We need to prove that Q' can also be obtained from Q by a sequence of mutations at sinks. But [Lampe, proof of Proposition 2.2.8] shows that Q' can be obtained from Q by a sequence of mutations at sinks and sources. Since every

To summarize, the outcome of our two-step process is that every arrow e of Q which starts at a vertex in A and ends at a vertex in B reverses its orientation, while all other arrows preserve their orientation. In other words, the outcome of our two-step process is the same as the outcome of turning all arrows of Q which start at a vertex in A and end at a vertex in B. But the latter outcome is  $\operatorname{mut}_{A,B} Q$  (because this is how  $\operatorname{mut}_{A,B} Q$  was defined), while the former outcome is  $\mu_b \left( \operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q \right)$  can be obtained from Q by our two-step process). Thus, we have obtained  $\mu_b \left( \operatorname{mut}_{A \cup \{b\}, B \setminus \{b\}} Q \right) = \operatorname{mut}_{A,B} Q$ . This proves (2).

<sup>9</sup>*Proof.* Assume the contrary. Then, there exists an arrow of Q which starts at a vertex in B and ends at a vertex in A. Let e be such an arrow. Then, e ends at a vertex in A. In other words,  $t(e) \in A = \{i\}$ , so that t(e) = i. In other words, e ends at e. But this is impossible, since e is a source. This contradiction proves that our assumption was wrong, qed.

<sup>10</sup> Proof. Each arrow of Q which starts at a vertex in A and ends at a vertex in B must be an arrow starting at i (because it starts at a vertex in  $A = \{i\}$ , but the only vertex in  $\{i\}$  is i). It thus remains to prove the converse − i.e., to prove that each arrow of Q starting at i is an arrow of Q which starts at a vertex in A and ends at a vertex in B. So let e be an arrow of Q starting at i. Then, e clearly starts at a vertex in A (since  $i \in \{i\} = A$ ). It remains to prove that e ends at a vertex in B. But Q is acyclic, and thus we cannot have s(e) = t(e) (since otherwise, the arrow e would form a trivial cycle). Hence,  $s(e) \neq t(e)$ . But s(e) = i (since e starts at i), so that  $t(e) \neq s(e) = i$  and thus  $t(e) \in Q_0 \setminus \{i\} = Q_0 \setminus A = B$ . Hence, e ends at a vertex in e. This

completes our proof.

 $t(e) \neq b$  (because if we had t(e) = b, then e would contradict (1)). But also  $t(e) \notin A$  (since e does not end at a vertex in A), so that  $t(e) \in Q_0 \setminus A = B$  and thus  $t(e) \in B \setminus \{b\}$  (since  $t(e) \neq b$ ). Hence, the arrow e ends at a vertex in  $B \setminus \{b\}$ . It also starts at a vertex in B; thus, it either starts at e or it starts at a vertex in e0. Our claim now follows from our analysis of Case 4 (in the case when e1 starts at e2 and from our analysis of Case 5 (in the case when e2 starts at a vertex in e3. In either case, our claim is proven.)

mutation at a source can be simulated by a sequence of mutations at sinks (by Exercise 0.1 **(b)**), this yields that Q' can be obtained from Q by a sequence of mutations at sinks. This solves Exercise 0.1 **(c)**.

## References

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